Reduced Complexity Interpolation for List Decoding Hermitian Codes

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Abstract—List decoding Hermitian codes using the Guruswami-Sudan (GS) algorithm can correct errors beyond half the designed minimum distance. It consists of two processes: interpolation and factorisation. By first defining a Hermitian curve, these processes can be implemented with an iterative polynomial construction algorithm and a recursive coefficient search algorithm respectively. To improve the efficiency of list decoding Hermitian codes, this paper presents two contributions to reduce the interpolation complexity. First, in order to simplify the calculation of a polynomial’s zero condition during the iterative interpolation, we propose an algorithm to determine the corresponding coefficients between the pole basis monomials and zero basis functions of a Hermitian curve. Second, we propose a modified complexity reducing interpolation algorithm. This scheme identifies any unnecessary polynomials during iterations and eliminates them to improve the interpolation efficiency. Due to the above complexity reducing modifications, list decoding long Hermitian codes with higher interpolation multiplicity becomes feasible. This paper shows list decoding algorithm can achieve significant coding gain over the conventional unique decoding algorithm.

Index Terms—List decoding, hermitian codes, decoding efficiency.

I. INTRODUCTION

LIST decoding of algebraic-geometric (AG) codes can outperform conventional unique decoding algorithms by correcting errors beyond $(d^*-1)/2$, where $d^*$ is the designed minimum distance of the AG code. For a $(n, k)$ AG code with length $n$ and dimension $k$, $d^* = n - k - g + 1$ where $g$ is the genus of the corresponding algebraic curve. Sudan [1] introduced the first list decoding algorithm for low rate Reed-Solomon (RS) codes, followed by Shokrollahi and Wasserman [2] to list decode low rate AG codes. Guruswami and Sudan [3, 4] later presented a complete version of a list decoding algorithm for all rate RS and AG codes, called the Guruswami-Sudan (GS) algorithm. For a $(n, k)$ AG code, the GS algorithm can correct up to

$$
\tau_{GS} = n - \left\lfloor \sqrt{n(n-d^*)} \right\rfloor - 1
$$

errors. The algorithm consists of two processes: interpolation, which builds an interpolated polynomial based on the received information, and factorisation, which finds the transmitted message information based on the interpolated polynomial.

They can be implemented with an iterative polynomial construction algorithm [5-8] and a recursive coefficient search algorithm [9-12] respectively.

RS codes are widely used in industry, such as communication and storage systems. Compared with RS codes defined over the same Galois field (GF), AG codes have better error-correction capability because they have longer code word lengths resulting in a large designed minimum distance [13, 14]. Therefore, AG codes could be considered as the successor of RS codes. Among them, Hermitian codes are one of the best performing and well developed AG codes. Johnston and Carrasco [15, 16] evaluated Hermitian code performance by employing the Sakata algorithm [17, 18] with majority voting [19] and inverse discrete Fourier transforms. This is a unique decoding algorithm with the half distance error-correction bound. To decode Hermitian codes beyond this bound, Hoholdt and Nielsen [5, 6] presented a mathematical framework of list decoding Hermitian codes with the GS algorithm. After an extensive study of [5, 6], the authors published the first simulation results of list decoding Hermitian codes in [20], which showed significant coding gains can be achieved over a unique decoding algorithm [17-19]. However, list decoding Hermitian codes with the GS algorithm remains complex and limits the application of the GS algorithm to longer codes. According to the complexity analysis in [8], the GS algorithm’s high complexity is mainly caused by the iterative interpolation, in which a set of polynomials is tested for different zero conditions and modified interactively. For Hermitian codes, to define the zero condition of a polynomial, we need to transfer it into a polynomial written with respect to the zero basis functions of a Hermitian curve, which is not very efficient for implementation. However, the zero condition of a polynomial can also be defined without this transfer based on knowledge of the corresponding coefficients between the pole basis monomials and zero basis functions of a Hermitian curve. Inspired by this, we propose an efficient algorithm to determine these coefficients before interpolation so that they can be applied afterwards. In order to improve the list decoding efficiency for RS codes, the authors [8] proposed a complexity reducing scheme to identify any unnecessary polynomials in the set and eliminate them during the iterative interpolation process. This scheme is also valid for list decoding Hermitian codes. This paper proposes a complexity reducing interpolation algorithm applying the scheme in [8]. The factorisation process of AG codes can be implemented by applying the recursive coefficient search algorithm [11]. A more generalised factorisation algorithm which can be applied to both RS and AG codes is later presented by the

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authors in [12]. By improving the efficiency of interpolation, list decoding longer Hermitian codes with higher multiplicity becomes feasible. This paper evaluates the performance of 4 different Hermitian codes defined in GF(16) and GF(64) over the additive white Gaussian noise (AWGN) and Rayleigh fading channels.

The structure of the paper is organised as follows: Section II presents prerequisite knowledge of Hermitian curves and Hermitian codes; Section III briefly describes the GS decoding of Hermitian codes, in which the zero condition of a polynomial and the decoding parameters are defined; Section IV proposes an efficient algorithm to determine the corresponding coefficients between the pole basis monomials and zero basis functions of a Hermitian curve; Section V presents the reduced complexity interpolation algorithm; simulation results evaluating the performance of 4 different Hermitian codes are presented in Section VI and conclusions are given in Section VII.

II. PREREQUISITE KNOWLEDGE

A Hermitian curve defined over GF(q), where q is the size of Galois field, is given as:

$$H_w(x, y, z) = x^{w^2 + 1} + y^w z + yz^w$$

(2)

where \(w = \sqrt{q}\) and has a genus \(g = \frac{w(w-1)}{2}\) [14]. For simplicity, this paper assumes GF(q) is an extension field of GF(2). \(H_w(x, y, z)\) is a projective curve with three affine components \(H_w(1, y, z)\), \(H_w(x, 1, z)\) and \(H_w(x, y, 1)\). By selecting one affine component \(H_w(x, y, 1)\), there are \(n = w^3\) affine points on it, which are of the form \(p_i = (x_i, y_i, 1)\) \(0 \leq i \leq n - 1\), and a point of infinity \(p_\infty = (0, 1, 0)\). Rational functions on the curve \(H_w\) must have a pole only at \(p_\infty\), and can be generally written as [14]:

$$\frac{x^\lambda y^\gamma z^\delta}{z^{\lambda + \gamma + \delta}} (0 \leq \lambda \leq w, \delta \geq 0)$$

(3)

where \(\lambda, \delta \in N\) and \(N\) denotes the set of nonnegative integers in this paper. Based on \(H_w(x, y, z) = 0\), \(\frac{z}{y} = \frac{x^{w^2 + 1} + y^w z + yz^w}{z^w x^w}\). Therefore, the pole orders [21] of variable \(\frac{z}{y}\) at \(p_\infty\) are:

$$v_{p_\infty} \left( \left( \frac{x}{z} \right)^{-1} \right) = w, v_{p_\infty} \left( \left( \frac{y}{z} \right)^{-1} \right) = w + 1.$$  

(4)

Pole basis \(L_w\) of curve \(H_w\) contains a set of rational functions with \(z = 1\), coefficients 1, and increasing pole orders. As \(z = 1\), these rational functions can be reduced to bivariate monomials, defined as [5]:

$$L_w = \{ \phi_a \mid v_{p_\infty} (\phi_a^{-1}) < v_{p_\infty} (\phi_{a+1}^{-1}), a \in N \}$$

(5)

where pole basis monomial \(\phi_a\) and its pole order can be generally written as:

$$\phi_a = x^\lambda y^\gamma z^\delta, v_{p_\infty} (x^\lambda y^\gamma z^\delta)^{-1} = w \lambda + (w + 1) \delta, 0 \leq \lambda \leq w, \delta \geq 0.$$  

(6)

Also, affine points can be distinctively denoted by their \(x, y\) coordinates as: \(p_i = (x_i, y_i)\). Nonnegative integers can be grouped into nongaps which are the pole orders of monomials in \(L_w\), and gaps otherwise. With respect to every affine point \(p_i\), there exists a zero basis \(Z_{w,p_i}\) which contains a set of rational functions \(\psi_{p_i,\alpha}\) with increasing zero orders [21] at \(p_i(z_{p_i})\), defined as [5]:

$$Z_{w,p_i} = \{ \psi_{p_i,\alpha} \mid v_{p_i} (\psi_{p_i,\alpha}) < v_{p_i} (\psi_{p_i,\alpha+1}), \alpha \in N \}.$$  

(7)

\(\psi_{p_i,\alpha}\) has a zero order \(\alpha\) at affine point \(p_i\) and it can be generally written as [6]:

$$\psi_{p_i,\alpha} = \psi_{p_i,\lambda+\delta} = (x - x_i)^\lambda (y - y_i)^\gamma z^\delta,$$

$$0 \leq \lambda \leq w, \delta \geq 0.$$  

(8)

The relationship between pole basis monomial \(\phi_a\) and zero basis function \(\psi_{p_i,\alpha}\) can be written as [6]:

$$\phi_a = \sum_{\alpha \in N} \gamma_{a,p_i,\alpha} \psi_{p_i,\alpha}$$  

(9)

where \(\gamma_{a,p_i,\alpha} \in GF(q)\) are the corresponding coefficients.

The construction of a \((n, k)\) Hermitian code can be described as evaluating the \(n\) affine points of \(H_w\) over the message polynomial \(f\):

$$T(n, k) = (c_0, c_1, \ldots, c_n-1) = (f(p_0), f(p_1), \ldots, f(p_{n-1}))$$

(10)

where \(c_0, c_1, \ldots, c_n-1 \in GF(q)\) and \(f\) could be written as:

$$f(x, y) = f_0\phi_0 + f_1\phi_1 + \ldots + f_{k-1}\phi_{k-1}$$  

(11)

and \(f_0, f_1, \ldots, f_{k-1} \in GF(q)\) are message symbols and \(\phi_0, \phi_1, \ldots, \phi_{k-1}\) are the first \(k\) monomials in \(L_w\).

To decode a \((n, k)\) Hermitian code with the GS algorithm, the pole order of variable \(z\) is defined as \(v_{p_\infty}(z^{-1}) = v_{p_\infty}(\phi_{k-1}^{-1})\). Then, the weighted degree of any trivariate monomial \(\phi_a z^b\) can be defined as:

$$deg_{w}(\phi_a z^b) = v_{p_\infty}(\phi_a^{-1}) + bv_{p_\infty}(z^{-1})$$  

(12)

and a lexicographic order (ord) can be defined to arrange the monomials \(\phi_a z^b\):

$$\phi_a z^{b_1} < \phi_a z^{b_2}$$

if \(deg_w(\phi_a z^{b_1}) < deg_w(\phi_a z^{b_2})\), or \(deg_w(\phi_a z^{b_1}) = deg_w(\phi_a z^{b_2})\) and \(b_1 < b_2\) [5]. \(F_q^w[x, y, z]\) is the ring of polynomials defined over the set of pole basis functions of the Hermitian curve \(L_w\), which can be generally written as:

$$f(x, y, z) = \sum_{a,b \in N} f_{ab}\phi_a(x) z^b,$$

where \(f_{ab} \in GF(q)\) and \(\phi_a \in L_w\). Subsequently, \(F_q^w[x, y]\) is a subset of \(F_q^w[x, y, z]\) with \(z\) degree equal to zero and \(v_{p_\infty}(\phi_a^{-1}) \leq u\). If we denote \(u_z = v_{p_\infty}(z^{-1})\), the message polynomial in (11) is a polynomial in \(F_q^w[x, y]\). The following definition is given for polynomials defined in \(F_q^w[x, y]\):

**Definition 1:** If \(\phi_a z^{b'}\) is the maximal monomial in polynomial \(f \in F_q^w[x, y, z]\):

$$\phi_a z^{b'} = \max \{ \phi_a z^b \mid f_{ab} \neq 0 \}.$$  

then \(\phi_a z^{b'}\) is called the leading monomial of \(f\) and its coefficient \(f_{ab}\) is called the leading coefficient of \(f\), denoted as:

$$LM(f) = \phi_a z^{b'}, \text{ and } LC(f) = f_{ab}$$  

The weighted degree of \(f\) \(deg_w(f)\) and leading order of \(f\) \((ord(f))\) are defined as:
\[ \deg_w(f) = \deg_w(\phi_a z^{b'}) \text{, and } \text{lod}(f) = \text{ord}(\phi_a z^{b'}) \]

Based on the above definitions, for any two polynomials \( f \) and \( h \in F_q[x, y, z] \), \( f < h \) if \( \text{lod}(f) < \text{lod}(h) \).

III. GS Decoding of Hermitian Codes

The GS algorithm consists of two processes: interpolation and factorisation. Given a received word \( R = (r_0, r_1, \ldots, r_{n-1}) \in \text{GF}(q) \), \( n \) interpolated units can be formed by combining each received symbol with its respective affine point used in encoding as: \((p_i, r_0), (p_i, r_1), \ldots, (p_i, r_{n-1})\).

Interpolation generates the minimal polynomial \( Q \in F_q[x, y, z] \) which has a zero of multiplicity of at least \( m \) over the received word. In general,

\[
Q = \sum_{a,b \in N} Q_{ab} \phi_a z^b
\]

where \( Q_{ab} \in \text{GF}(q) \). For each interpolated unit \((p_i, r_1)\), the polynomial \( Q \) can also be written with respect to the zero basis functions in \( \Phi_{a,p_i} \) as [6]:

\[
Q = \sum_{\alpha, \beta \in N} Q_{\alpha \beta} p_{\alpha \beta} \phi_{\alpha} (z - r_i)^{\beta}
\]

where \( Q_{\alpha \beta} \in \text{GF}(q) \). Based on equation (8), clearly, \( Q(p_i, r_1) = 0 \). Furthermore, if \( Q_{\alpha \beta}(p_i, r_1) = 0 \) for \( \alpha + \beta < m \), polynomial \( Q \) has a zero of multiplicity of at least \( m \) at unit \((p_i, r_1)\) [5, 6]. As \( z^b = (z - r_i)^{r_i} \sum_{\beta \leq b} \binom{b}{\beta} r_i^{\beta} (z - r_i)^\beta \) and \( \phi_a = \sum_{\alpha \in N} \gamma_{a,p_i} \phi_{\alpha} \), substituting them into (13) gives:

\[
Q = \sum_{a,b \in N} Q_{ab} \sum_{\alpha \in N} \gamma_{a,p_i} \phi_{\alpha} (z - r_i)^{\beta}
\]

Therefore, coefficients \( Q_{\alpha \beta} \) of (14) can be written as:

\[
Q_{\alpha \beta} = \sum_{a,b \geq \beta} Q_{ab} \binom{b}{\beta} \gamma_{a,p_i} \phi_{\alpha} (z - r_i)^{\beta}.
\]

(16) defines the zero condition constraints of the coefficients \( Q_{ab} \) of polynomial \( Q \), so that \( Q \) has a zero of multiplicity of at least \( m \) over unit \((p_i, r_i)\). If we denote the constraints in (16) to the coefficients of polynomial \( Q \) as \( D_{\alpha \beta}^{(p_i, r_i)}(Q) \), such that:

\[
D_{\alpha \beta}^{(p_i, r_i)}(Q) = Q_{\alpha \beta} = \sum_{a,b \geq \beta} Q_{ab} \binom{b}{\beta} \gamma_{a,p_i} \phi_{\alpha} (z - r_i)^{\beta}.
\]

(17) Then after interpolating all the affine points, the interpolated polynomial \( Q \) shall satisfy:

\[
Q = \min_{\text{lod}(Q)} \{ Q \in F_q[x, y, z] \mid D_{\alpha \beta}^{(p_i, r_i)}(Q) = 0, \text{ for } i = 0, 1, \ldots, n - 1 \wedge \alpha + \beta < m \}.
\]

As there are \( \binom{m+1}{2} \) permutations of \((\alpha, \beta)\) for \( \alpha + \beta < m \), there are in total:

\[
C = n \binom{m+1}{2}
\]

zero condition constraints that the coefficients \( Q_{ab} \) of polynomial \( Q \) need to satisfy. \( C \) also represents the number of iterations in the interpolation algorithm [5, 6], in which each iteration imposes a zero condition constraint on \( Q_{ab} \). The weighted degree upper bound of interpolated polynomial \( Q \) is defined as [5, 6]:

\[
\max \{ \deg_w Q \} = l_m v_p(\alpha^{-1}) + t_m
\]

where \( l_m \) is the maximal number of output candidates from factorisation, defined as:

\[
l_m = \max \{ u \mid \frac{u}{2} v_p(\alpha^{-1}) - (u - 1)g \leq C \} - 1
\]

and parameter \( t_m \) is defined as:

\[
t_m = \max \{ u \mid (l_m + 1)u - \Gamma(u) + \left( \frac{lm + 1}{2} \right) v_p(\alpha^{-1}) - l_m g \leq C \}
\]

where \( u \in N \) and \( \Gamma(u) \) denotes the number of gaps that are less than or equal to the nonnegative integer \( u \).[5]

According to Theorem 3 of [5], if there exists a polynomial \( h \in \Phi_{a,p_i} \) such that:

\[
\Lambda(h, R) > \deg_w Q
\]

where \( \Lambda(h, R) = \{ t \mid h(p_i) = r_i, i = 0, 1, \ldots, n - 1 \} \) represents the number of affine points that satisfy \( h(p_i) = r_i \), then \( h \) is the \( z \) root of \( Q \), i.e. \( Q(x, y, h) = 0 \), or equivalently \((z - h) \mid Q(x, y, z)\). Factorisation finds the \( z \) roots of the interpolated polynomial \( Q \), among which the message polynomial of (11) is included [10-12, 20]. Therefore, the GS algorithm’s error-correction capability \( \tau_m \) is:

\[
\tau_m = n - \Lambda(h, R) = n - \left[ \frac{\deg_w Q}{m} \right] - 1.
\]

As the upper bound of \( \deg_w Q \) is defined by (20), then:

\[
\tau_m \geq n - \left[ \frac{l_m v_p(\alpha^{-1}) + t_m}{m} \right] - 1.
\]

IV. Determining The Corresponding Coefficients

Based on (17), the corresponding coefficients \( \gamma_{a,p_i} \) are critical for defining the zero condition of a polynomial in \( F_q[x, y, z] \). Without the knowledge of them, we must transfer a general polynomial (13) into (14) and find the coefficients \( Q_{\alpha \beta} \), which is not efficient during the iterative interpolation. In fact, the corresponding coefficients \( \gamma_{a,p_i} \) can be determined independently of the received word. Therefore, if they can be determined beforehand and applied during the iterations, the interpolation efficiency can be greatly improved.

This section proposes an algorithm to determine them.

The problem we intend to solve can be simply stated as: given an affine point \( p_i = (x_i, y_i) \) of the curve \( H_w \) and a pole basis monomial \( \phi_a \), determine the corresponding coefficients \( \gamma_{a,p_i} \) so that \( \phi_a \) can be written as a sum of the zero basis functions \( \psi_{p_i, \alpha} \): \( \phi_a = \sum_{\alpha \beta \in N} \gamma_{a,p_i} \psi_{p_i, \alpha} \). For any two pole basis monomials \( \phi_a \) and \( \phi_{a_2} \) in \( L_w \), \( \phi_a \phi_{a_2} = \sum_{\alpha \in N} \phi_a \psi_{p_i, \alpha} \) and the zero basis function \( \psi_{p_i, \alpha} \) in (8) can be written as a sum of pole basis monomials \( \phi_a \) as [6]:

\[
\psi_{p_i, \alpha} = \sum_{a \in N} \phi_a
\]
where coefficients $\zeta \in GF(q)$. Based on (8), we partition 
$\psi_{p, \alpha}$ as:
$$\psi_{p, \alpha} = \psi_{p, \alpha}^A \psi_{p, \alpha}^B$$
where $\psi_{p, \alpha}^A = (x - x_i)^{3}$ and $\psi_{p, \alpha}^B = (y - y_i) - x_i (x - x_i)^{3}$. It is easy to recognise that $\psi_{p, \alpha}^A$ has leading monomial $LM(\psi_{p, \alpha}^A) = x^3$ and leading coefficient $LC(\psi_{p, \alpha}^A) = 1$. Based on (27), $\psi_{p, \alpha}^{B}$ has leading monomial $LM(\psi_{p, \alpha}^B) = y^\delta$ and leading coefficient $LC(\psi_{p, \alpha}^B) = 1$. Therefore in the new accumulated polynomial resulting in $\phi_a$. While $\tilde{\psi} \neq \phi_a$, in (31), the second largest monomial $\phi_{L-1}$ with coefficient $\zeta_{L-1}$ is identified by step 5. Then, find another zero basis function $\psi_{p, \alpha}$ in $Z_{w,p}$, where $LM(\psi_{p, \alpha}) = \phi_{L-1}$. According to Corollary 1, this zero basis function always exists and it can be written as: $\psi_{p, \alpha} = \sum a' \in N, a' < L - 1 \cdot \zeta_a \phi_{a'} + L_{\alpha}$. At step 6, the corresponding coefficient between monomial $\phi_a$ and the found zero basis function $\psi_{p, \alpha}$ can be determined as: $\gamma_{a, p, \alpha} = \zeta_{L-1}$. As a result, the accumulated calculation of step 7 can be written as:
$$\tilde{\psi} = \sum a' \in N, a' < L - 1 \cdot \zeta_a \phi_{a'} + L_{\alpha} + \sum a' \in N, a' < L - 1 \cdot \zeta_a \phi_{a'} + L_{\alpha} - \zeta_{L-1} \zeta_{a'} \phi_{a'} - \phi_{L-1} = \sum a' \in N, a' < L - 1 \cdot \zeta_a \phi_{a'} + \phi_a.$$ 

Therefore in the new accumulated $\tilde{\psi}$, $\zeta_{L-1} \phi_{L-1}$ is eliminated while the leading monomial $\phi_a$ is preserved. If the updated $\tilde{\psi} \neq \phi_a$, its second largest monomial $\phi_{L-1}$ is again eliminated while $\phi_a$ is always preserved as a leading monomial by the same process. The algorithm terminates after all monomials that are smaller than $\phi_a$ have been eliminated and results in $\tilde{\psi} = \phi_a$. This process is equivalent to the sum calculation of (9). Here a worked example is presented to illustrate Algorithm A.

Example: Given $p_1 = (\sigma^2, \sigma^2)$ ($\sigma$ is a primitive element in $GF(4)$ satisfying $\sigma^2 + \sigma + 1 = 0$) is an affine point on curve $H_2$ and a pole basis ($L_2$) monomial $\phi_5 = y^2$, determine the corresponding coefficients $\gamma_{5, p, \alpha}$ so that $\phi_5$ can be written as $\phi_5 = \sum a \in N \gamma_{5, p, \alpha} \psi_{p, \alpha}$.

Based on (8), the first 8 zero basis functions in $Z_{2,p}$, can be listed as:
$$\psi_{p, 0} = (x - x_0)^2 = 1;$$
$$\psi_{p, 1} = (x - x_1)^2 = \sigma^2 + x;$$
$$\psi_{p, 2} = (x - x_2)^2 = \sigma + x^2;$$
$$\psi_{p, 3} = (y - x_3)^2 = \sigma (x - x_3) = \sigma + sy$$
$$\psi_{p, 4} = (x - x_4)^2 = \sigma (x - x_4)^2 = 1 + \sigma x^2 + \sigma^2 y^2 + \sigma x^2 + x y^2;$$
$$\psi_{p, 5} = (x - x_5)^2 = \sigma (x - x_5)^2 = \sigma^2 + \sigma x^2 + \sigma x^2 + \sigma x^2 + x y^2;$$
$$\psi_{p, 6} = y^2 = \sigma^2 + \sigma x^2 + x y^2;$$
$$\psi_{p, 7} = y^2 = \sigma^2 + \sigma x^2 + x y^2.$$

 Initialise all $\gamma_{5, p, \alpha} = 0$. In $Z_{2,p}$, as $LM(\psi_{p, 0}) = \phi_5$, we let $\gamma_{5, p, 0} = 1$ and initialise the accumulated polynomial $\psi = \psi_{p, 0} = \sigma^2 + \sigma x^2 + x^2$.

As $\psi \neq \phi_5$, its second largest monomial $\phi_{L-1} = x^2$ with coefficient $\zeta_{L-1} = \sigma^2$ is identified. Among the zero basis functions in $Z_{2,p}$, we find $\psi_{p, 2}$ with $LM(\psi_{p, 2}) = \phi_{L-1} = x^2$ and let $\gamma_{5, p, 2} = \zeta_{L-1} = \sigma^2$. Update $\psi = \psi + \gamma_{5, p, 2} \psi_{p, 2} = \sigma + y^2$.

As $\psi \neq \phi_5$, again its second largest monomial $\phi_{L-1} = 1$ with coefficient $\zeta_{L-1} = 1$ are identified. Among the zero basis functions in $Z_{2,p}$, we find $\psi_{p, 2} = \psi_{p, 0} = LM(\psi_{p, 0}) = L_{\alpha} = 1$. The algorithm terminates after all monomials that are smaller than $\phi_a$ have been eliminated and results in $\tilde{\psi} = \phi_a$. This process is equivalent to the sum calculation of (9). Here a worked example is presented to illustrate Algorithm A.

Proof: Notice that functions $\psi_{p, \alpha}$ with $LM(\psi_{p, \alpha}) > \phi_a$ will not contribute to the sum calculation of (9) and their corresponding coefficients $\gamma_{a, p, \alpha} = 0$. The zero basis function $\psi_{p, \alpha}$ found at step 2 has leading monomial $\phi_L = \phi_a$. Based on (30), it can be written as:
$$\psi_{p, \alpha} = \sum a' \in N, a' < L - 1 \cdot \zeta_a \phi_{a'} + L_{\alpha}.$$ 

(31) indicates the corresponding coefficient between $\phi_a$ and $\psi_{p, \alpha}$ is: $\gamma_{a, p, \alpha} = 1$. Polynomial $\tilde{\psi}$ initialised by step 3 is
1, and let $\gamma_{5,p_1,0} = \zeta_{L-1} = \sigma$. Update $\hat{\psi} = \psi + \gamma_{5,p_1,0}\psi_{p_1,0} = y^2$.

Now, $\hat{\psi} = \phi_5$, we can stop the algorithm and output $\gamma_{5,p_1,0} = \sigma$, $\gamma_{5,p_1,2} = \sigma^2$ and $\gamma_{5,p_1,6} = 1$. The rest of the corresponding coefficients $\gamma_{5,p_1,\alpha}$ for $\alpha \neq 0, 2, 6$.

Before interpolation, the monomials $\phi_5$ that exist in the interpolated polynomial $Q$ are unknown. However, the weighted degree upper bound of polynomial $Q$ is defined by (20), from which the largest pole basis monomial $\phi_{5,max}$ that might exist in $Q$ can be predicted by $v_{p_{max}}(\phi_{5,max}) = \max \{ \deg_u Q \}$. Based on an interpolation multiplicity $m$, with parameter $\alpha < m$, the corresponding coefficients that might be used in interpolation are $\gamma_{0,p_1,\alpha} \sim \gamma_{5,max,p_\alpha}$. Therefore Algorithm $A$ can be used to determine all the corresponding coefficients $\gamma_{0,p_1,\alpha} \sim \gamma_{5,max,p_\alpha}$ and only $\gamma_{0,p_1,\alpha} \sim \gamma_{5,max,p_\alpha}$ for $\alpha < m$ are stored for interpolation in order to minimise the memory requirement. For example, to list decode the $(8, 4, 4)$ Hermitian code with multiplicity $m = 2$, $\max \{ \deg_u Q \} = 13$. Therefore, the largest pole basis monomial that might exist in $Q$ is $\phi_{max} = \phi_{12} = x^2 y^5$ and Algorithm $A$ can be applied to calculate all the corresponding coefficients $\gamma_{0,p_1,\alpha} \sim \gamma_{12,p_1,\alpha}$ and $\gamma_{0,p_1,\alpha} \sim \gamma_{12,p_1,\alpha}$ for $\alpha < 2$ are stored.

V. Complexity Reducing Interpolation

Interpolation determines the polynomial $Q$ defined by (18). This can be implemented using an iterative polynomial construction algorithm [5-8]. First, a set of polynomials is initialised. During the iterations, they are tested by different zero conditions and modified interactively. As mentioned in Section III, there are in total $C$ iterations as given in (19), after which the minimal polynomial in the set is chosen as the interpolated polynomial $Q$. According to the iterative process analysis [8], the interpolated polynomial $Q$ has leading order $\text{lod}(Q) \leq C$. This indicates that those polynomials with leading order over $C$ will not be the chosen candidates. Also, if there is a polynomial in the set with leading order over $C$ during the iterations, the chosen polynomial $Q$ has not been modified with this polynomial, otherwise $\text{lod}(Q) > C$ [8]. Therefore, those polynomials with leading order greater than $C$ can be eliminated from the set during iterations in order to reduce unnecessary computations.

If $f \in F_7[x, y, z]$ has leading monomial $LM(f) = \phi_5 z^5$, polynomials in $F_7[x, y, z]$ can be partitioned into the following classes according to the $z$ degree of their leading monomial and the pole order $v_{p_{max}}(\phi_5^{-1})$ of $\phi_5$ as:

$$V_{\lambda+w} = \{ f \in F_7[x, y, z] \mid b' = \delta \wedge v_{p_{max}}(\phi_5^{-1}) = uw + \lambda, \quad \text{LM}(f) = \phi_5 z^{b'}, (\delta, u, \lambda) \in \mathbb{N}, \lambda < w \}$$

such that $F_7[x, y, z] = \bigcup_{\lambda, \delta \in \mathbb{N}, \lambda < w} V_{\lambda+w}$. According to Section III, the factorisation outputs are the $z$ roots of $Q$. Therefore, the $z$ degree of $Q$ is less than or equal to the maximal length of the output list $l_m$ of (21) and $Q$ is a polynomial chosen from the following classes:

$$V_j = V_{\lambda+w}, \quad (0 \leq \lambda < w, 0 \leq \delta \leq l_m).$$

At the beginning of the iterative process, a set of polynomials is initialised to represent each of the polynomial classes defined by (34) as:

$$G = \{ Q_j = Q_{\lambda+w} = y^\lambda z^\delta, Q_j \in V_j \}. \quad (35)$$

During the iterations, each polynomial $Q_j$ in the set $G$ is the minimal polynomial within its class $V_j$ that satisfies all the tested zero conditions. At the beginning of each iteration, the polynomial set $G$ is modified by:

$$G = \{ Q_j \mid \text{lod}(Q_j) \leq C \}$$

in order to eliminate those polynomials with leading order over $C$. Then the remaining polynomials in $G$ are tested by the zero condition defined in (17) as:

$$\Delta_j = D_{\alpha\beta}(Q_j). \quad (36)$$

The determined corresponding coefficients $\gamma_{a,p_i,\alpha}$ are applied for this calculation. Those polynomials with $\Delta_j = 0$ satisfy the zero condition and do not need to be modified. However, those polynomials with $\Delta_j \neq 0$ need to be modified. Among them, find the index of the minimal polynomial as $j'$ and record the minimal polynomial as $Q'$:

$$j' = \text{index} \{ \min \{ Q_j \mid \text{lod}(Q_j) \leq C \} \}. \quad (38)$$

$$Q' = Q_{j'}. \quad (39)$$

For $Q_{j'}$, it is modified as:

$$Q_{j'} = (x - x_i)Q' \quad (40)$$

where $x_i$ is the $x$ coordinate of affine point $p_i$ in the current interpolated unit $(p_i, v_i)$. The modified $Q_j$ satisfies $D_{\alpha\beta}(Q_{j'}) = 0$ because $D_{\alpha\beta}(Q_{j'}) = D_{\alpha\beta}(Q_j - \beta x_i) = D_{\alpha\beta}(Q_j)(x - x_i) = x_i \Delta_j - x_i \Delta_j' = 0$. The rest of the polynomials with $\Delta_j \neq 0$ are modified as:

$$Q_j = \Delta_j Q_j - \Delta_j Q_j'. \quad (41)$$

The modified $Q_j$ satisfies $D_{\alpha\beta}(Q_{j'}) = 0$ because $D_{\alpha\beta}(Q_{j'}) = \Delta_j Q_{j'} = \Delta_j Q_j - \Delta_j Q_j' = 0$. After $C$ iterations, the minimal polynomial in the set $G$ is chosen as the interpolated polynomial $Q$:

$$Q = \min_{\text{lod}(Q_j)} \{ Q_j \mid Q_j \in G \}. \quad (42)$$

From the above description, it can be seen that by applying the complexity reducing scheme in (36), the zero condition calculation in (37) and modifications in (40) and (41) for those polynomials $Q_j$ with $\text{lod}(Q_j) > C$ can be avoided, and therefore the interpolation efficiency can be improved. According to [8], this complexity reducing scheme is error dependent and can reduce complexity more significantly in low error weight situations. This is because the modification scheme in (36) takes action in earlier iteration steps for low error weight situations, and therefore computations can be reduced. Fig. 1 shows the interpolation (with different multiplicity $m$) complexity reduction by applying the scheme in (36) for decoding the $(64, 19, 40)$ Hermitian code. It is shown that complexity can be reduced significantly in low error weight situations, especially when $m = 1$, and complexity can be reduced by up to 48.83%. However, in high error...
weight situations, complexity reduction is not as significant. Based on Fig. 1, it can also be observed that the complexity reduction also depends on the interpolation multiplicity $m$. When $m = 1$, complexity reduction is significant; when $m = 2$, complexity reduction is marginal.

Summarising Section IV and V, the reduced complexity interpolation process for GS decoding Hermitian codes can be stated as:

**Initial computation:** Apply Algorithm A to determine all the necessary corresponding coefficients $\gamma_{\alpha, \beta, \alpha}$ and store them to be used by the iterative polynomial construction algorithm (Algorithm B).

Algorithm B: Iterative Polynomial Construction.

Initialisation: Initialise the set of polynomials $G$ by (35);

Step 1: For each interpolated unit $(p_i, r_i) (i = 0, 1, \cdots, n-1)$ {

Step 2: For each pair of the zero condition parameters $(\alpha, \beta)(\alpha + \beta < m)$ {

Step 3: Modify polynomial set $G$ by (36);

Step 4: Test the zero condition $\Delta_j$ of each polynomial in $G$ by (37);

Step 5: For polynomials $Q_j$ with $\Delta_j \neq 0$ {

Step 6: Denote the index of the minimal polynomial as $j'$ by (38) and record it as $Q'$ by (39);

Step 7: If $j = j'$, $Q_j$ is modified by (40);

Step 8: If $j \neq j'$, $Q_j$ is modified by (41);
}
}

At the end of the iterations, the minimal polynomial $Q$ is chosen from the set $G$ as (42).

**VI. PERFORMANCE EVALUATION**

Employing the above efficiency improved interpolation, list decoding of longer Hermitian codes with different multiplicity is feasible. The authors have implemented the GS decoding of Hermitian codes [3, 5] using C programming language, in which the factorisation process is implemented by the recursive coefficient search algorithm [11, 12]. The evaluating list decoder structure is presented in Fig. 2. Figs. 3 and 4 present the performance of (64, 19, 40) and (64, 39, 20) Hermitian codes, while Figs. 5 and 6 present the performance of (512, 153, 332) and (512, 289, 196) Hermitian codes. Simulations are run over AWGN and Rayleigh fading channels using the QPSK modulation. The Rayleigh fading channel is frequency nonselective with Doppler frequency $126.67$ Hz and date rate of $30$ kb/s. Over the fading channel, a block interleaver with size $100 \times n$ is used, where $n$ is the length of the code. Their performances are evaluated by measuring their coding gains (dB) over the unique decoding algorithm [15-19] at a bit error rate (BER) of $10^{-5}$.

According to the interpolation description in Section V, for interpolation with multiplicity $m$, there are $\omega(l_m + 1)$ polynomials being initialised that take part in $C$ iterations. Even though some of them will be eliminated during the iteration by the scheme in (36), a high value of $m$ will still lead to infeasibility for implementation as both the iteration number and polynomial set size grow with $m$ exponentially. Therefore, we have achieved some feasible results from the GS algorithm ($m = 1, 2$). Table I presents our simulation parameters for these 4 Hermitian codes. Before interpolation, Algorithm A is applied to determine the corresponding coefficients $\gamma_{\alpha, \beta, \alpha} \sim \gamma_{\max, \alpha, \alpha}$ and only $\gamma_{0, \alpha, \alpha} \sim \gamma_{\max, \alpha, \alpha}$ ($\alpha = 0, 1$) are stored in order to minimise the memory requirement. From Table I, it can be observed that to achieve the optimal result from the GS algorithm remains too complex for implementation. For example, to optimally decode (64, 19, 40) Hermitian code, there are 116 polynomials taking part in 9792 iterations in interpolation. But if assuming the GS algorithm is able to correct $\tau_{GS}$ errors and the transmitted code word $\bar{c}$ is known by the decoder, the theoretical optimal performance of the GS algorithm can also be evaluated without employing the interpolation and factorisation processes. This is achieved by measuring the Hamming distance between the received word $R$ and the transmitted code word $\bar{c}$ if it is not greater than $\tau_{GS}$. Decoding is claimed to be successful; otherwise, decoding is failure. Figs. 3 to 6 show that the GS algorithm approaches its optimal result with increasing interpolation multiplicity $m$. Among the performance evaluations, it is worth highlighting Fig. 4 which shows GS decoding the (64, 39, 20) Hermitian code with multiplicity $m = 2$ is close to the theoretical optimal result.

Table II analyses the simulation results shown in Figs. 3 to 6. During our simulations, the average number of errors $\tau_m$ that the GS algorithm is able to correct in order to achieve the corresponding performance is measured. Based on Table II, it can be observed that the GS algorithm’s coding gains grow with interpolation multiplicity $m$ and they are especially significant over the Rayleigh fading channel.

Note: indicates the pre-calculation step of the efficient list decoding algorithm.

**Fig. 2. Efficiency improved list decoder structure for Hermitian codes.**
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1.0E -06
1.0E -05
1.0E -04
1.0E -03
1.0E -02
1.0E -01
1.0E +00
-4 -3 -2 -1 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 BER

(a) over AWGN channel

(b) over Rayleigh fading channel

Fig. 3. List decoding performance of Hermitian code (64, 19, 40).

For example, GS decoding (64, 19, 40) Hermitian code with multiplicity \( m = 2 \) can achieve 1.42 dB coding gain over the Rayleigh fading channel. The GS algorithm can achieve more significant coding gains for low rate codes, but this is at the expense of higher decoding complexity. For example, comparing the (512, 153, 332) and (512, 289, 196) Hermitian codes, more significant coding gains can be achieved for the former. However, according to Table I, GS decoding (512, 153, 332) Hermitian code has higher decoding complexity because there are more polynomials being initialised that take part in the iterative interpolation. The same result can also be found by comparing the (64, 19, 40) and (64, 39, 20) Hermitian codes.

VII. CONCLUSION

This paper presents two contributions to reduce interpolation complexity so as to improve the efficiency of list decoding Hermitian codes. We first propose an efficient algorithm to determine the corresponding coefficients between the pole basis monomials and zero basis functions of a Hermitian curve. The coefficients are stored to be applied during the iterative interpolation in order to simplify the zero condition...
calculation of a polynomial. We then propose a complexity reducing interpolation algorithm by applying a scheme developed by the authors which eliminates any unnecessary polynomials during iterations. It is shown that this scheme can improve the interpolation efficiency significantly. Given that complexity can be reduced by the proposed methods, we have evaluated list decoding performance of longer Hermitian codes with higher interpolation multiplicity. Our results show that the GS algorithm can achieve significant coding gains over the unique decoding algorithm. The GS algorithm’s coding gains increase with interpolation multiplicity and it is more significant for low rate codes. However, this performance is at the expense of higher decoding complexity.

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